Social Learning in Multi Agent Multi Armed Bandits

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Joint Work with

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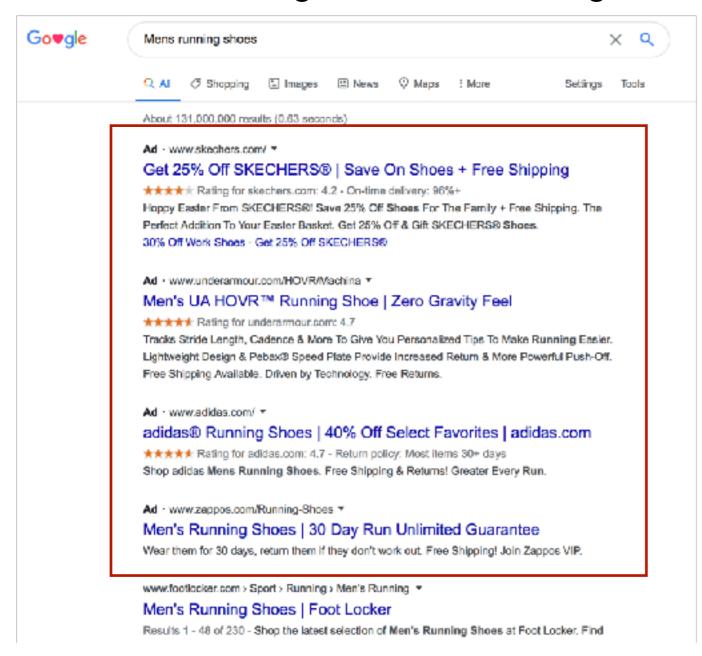
Drugs with a-priori unknown cure rates

Explore/Exploit Tradeoff for each new patient [Thompson' 33]

Exploit Prescribe a drug that has shown the best promise so farExplore Try a new drug to discover more promising alternativesRun a risk of not curing these patients

Explore/Exploit Phenomena

Online Advertising - Which of the large collection of ads to display?



Exploit - What has worked in the past?

Explore - Discover a more relevant ad

Internet Advertising accounts for 1 Trillion in revenue 6% of US GDP! [HBS Report 2019]

Outline

1. Single Agent MAB

2. The Multi-Agent Setup

3. Social Learning Algorithm

4. Insights

At each time, $t \in \{1, \dots, T\}$ an agent

- chooses an arm $I_t \in \{1, \cdots, K\}$
- receives a stochastic reward $X_t \in \{0,1\}$

$$\mathbb{P}[X_t = 1 | I_t] = \mu_{I_t}$$
 independent of everything else

Each arm corresponds to a drug in the previous example

Goal - Maximize total reward $\mathbb{E}[\sum_{t=1}^{T} X_t]$

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$$\mathbb{E}[\sum_{t=1}^{T} X_t]$$

<u>Challenge</u> Arm-means $(\mu_i)_{i=1}^K$ initially unknown

As we play arms, can learn $(\mu_i)_{i=1}^K$

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Explore-Exploit Tradeoff

Exploit Play the arm that has been best so far

Explore Play an arm played few times so as to see if it is good

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Performance Metric - Regret
$$R_T=\mu^*T-\mathbb{E}[\sum_{t=1}^{T}X_t]$$

$$\mu^*=\max\{\mu_1,\cdots,\mu_K\}$$

How much loss due to lack of knowledge?

Upper Confidence Bound (UCB) Algorithm

UCB Algorithm [Auer et.al. '02]

At time t, choose arm
$$I_t \in \arg\max_k \left(\widehat{\mu}_k(t-1) + \sqrt{\frac{4\alpha \log(t)}{N_k(t-1)}} \right)$$

$$\widehat{\mu}_k(t-1)$$
 Empirical Mean of arm k at time t-1

 $N_k(t-1)$ Number of times arm k has been played

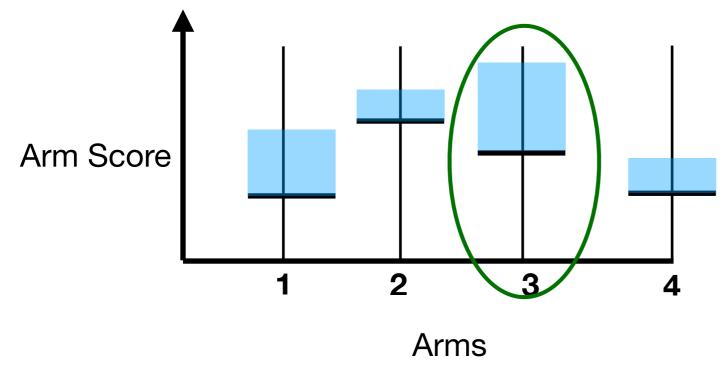
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Optimism in the Face of Uncertainty

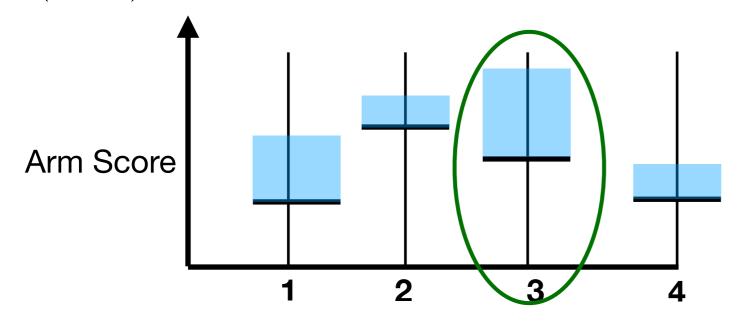
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Arms

Optimism in the Face of Uncertainty

Theorem
$$R_T \le O\left(\frac{K}{\Delta}\log(T)\right)$$

 Δ Difference in arm mean between best and second best arm

Multi Agent Setup



shorturl.at/huO57

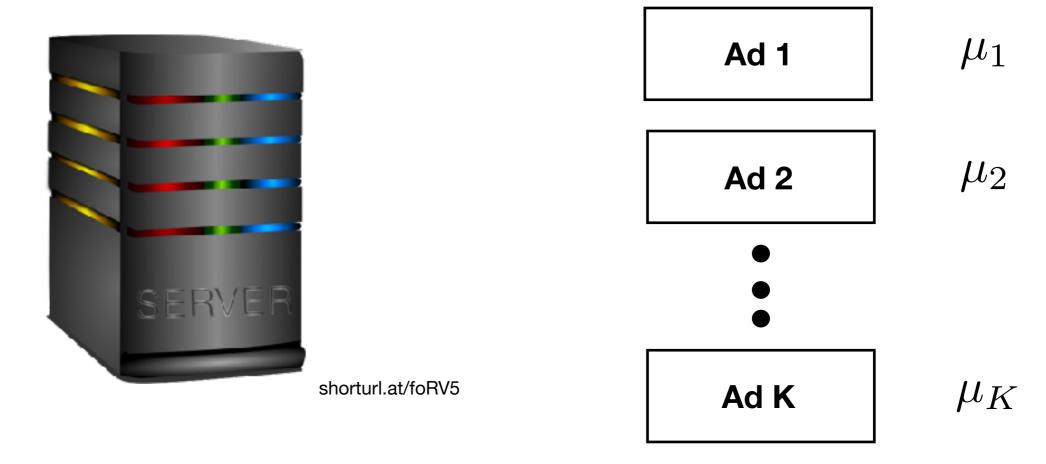
What if *multiple agents* play the *same* MAB instance ?

Can they collaborate and jointly reduce their individual regret?

One server is serving ads for a <u>fixed</u> keyword

At each search request, server can choose to display one ad

Choice of an arm to pull



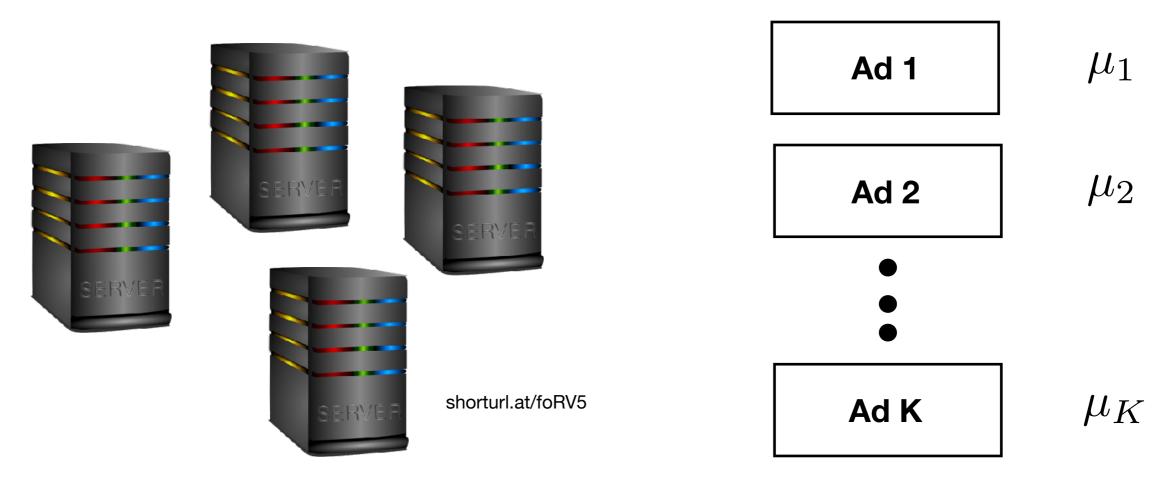
At the end, receives a stochastic reward

Goal is to maximize revenue (minimize regret)

Multiple servers serving ads for a fixed keyword

Each search request, routed to a server

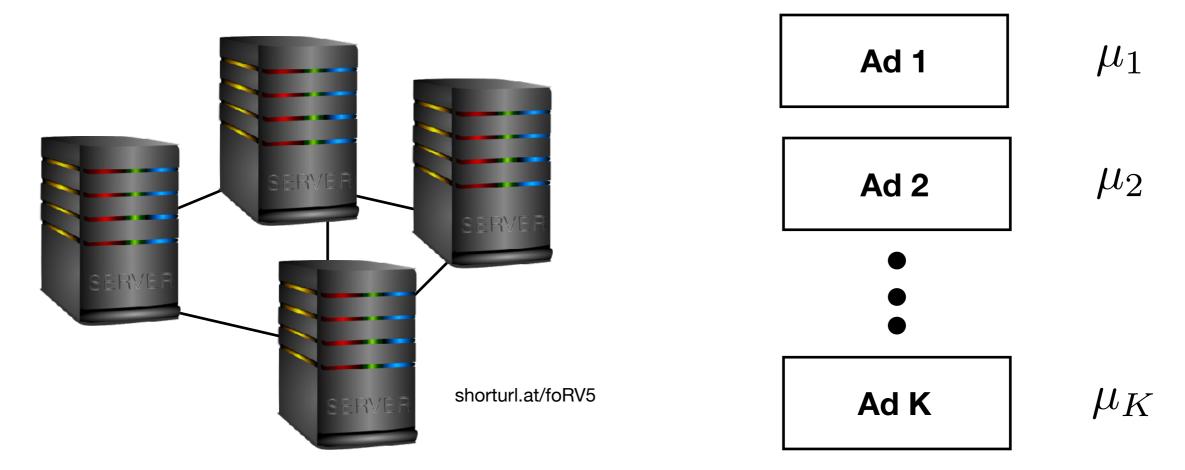
Each server chooses to display one ad when routed to it.



Multiple servers serving ads for a fixed keyword

Each search request, routed to a server

Each server chooses to display one ad when routed to it.



Servers can potentially collaborate and learn from each other's experience.

At each time, every server makes a decision from K alternatives

Large volume of search queries



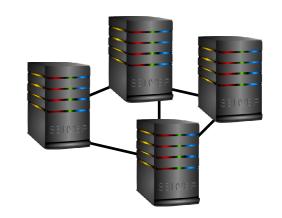
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Large volume of search queries

No Communication -

Individual Server Regret - $O\left(\frac{K}{\Delta}\log(T)\right)$

Communication Resources - 0



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1. No Communication -

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Individual Server Regret -
$$O\left(\frac{1}{N}.\frac{K}{\Delta}\log(T)\right)$$

Overall system can be abstracted as a single agent

Communication Resources - T broadcasts per agent!

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Communication Resources - T broadcasts per agent!

Best of both situations ??

The Multi Agent Problem

Karms, N agents,

Asynchronous System - Each agent j has an i.i.d. Poisson clock $\mathcal{C}_j(\cdot)$

When clock $\mathcal{C}_j(\cdot)$ rings for the t^{th} time, agent j

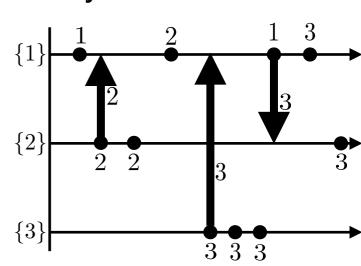
- 1. Plays an arm $I_j(t) \in \{1, \dots, K\}$ and receives reward $X_j(t) \in \{0, 1\}$
- 2. Can choose to send information to any other agent of choice

$$\mathbb{P}[X_j(t)=1|I_j(t)]=\mu_{I_j(t)}$$
 Independent rewards across agents

Decentralized Algorithms -

Choice of arm of an agent only a function of its observed history.

 $\underline{\mathsf{Minimize\ Individual\ Regret}} - R_T^{(j)} := \mu^*T - \mathbb{E}\left[\sum_{t=1}^T \mu_{I_j(t)}\right]$

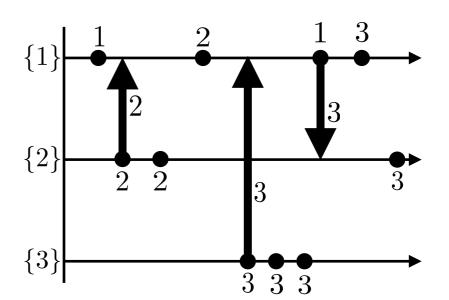


The Multi Agent Problem

Communication Constraints

1. Pairwise communication

Desirable for practical implementations



2. Fixed number of bits per message

Can't communicate arm-means to arbitrary precision.

Number of bits must not depend on arm-gap or other parameters of the instance.

3. Each agent communicates only o(T) times

In the search engine example, each server gets a huge volume of search queries.

If servers communicate as frequently as they make decisions, it will congest the server system

Related Work - Multi Agent Bandits

Multi-Agent Bandits are receiving wide attention -

1. Competitive Agents -

Multiple agents pull same arm, no-one (or only a subset) get rewards

[Anandkumar et.al '11],[Kalathil et.al. '14],[Rosenski et.al. '16],[Bistritz et.al. '18]

2. Collaborative Agents - Neighbors can observe all samples

[Buccapatnam et.al. '15][Landgren. et.al '16][Kolla et.al. '18] [Martinez-Rubio et.al '18]

Main Result

Theorem (Informal) - We give a social algorithm, where regret of any agent is

$$O\left(\left(\frac{K}{N} + \log(N)\right) \frac{\log(T)}{\Delta}\right) + O\left(\frac{\log^3(N)}{\Delta^2} \log\log(N)\right)$$
Constant Independent of time

- 1. Each agent communicates only O(log(T)) times
- 2. Each communication exchanges O(log(K)) bits (just arm-ids)
- 3. Each agent communicates with an agent chosen at random (Gossip)

Main Result

In order to derive insight, suppose K=N

	No-Interaction	Social Learning	Full Interaction
Regret	$O\left(N\frac{\log(T)}{\Delta}\right)$	$O\left(\log(N)\frac{\log(T)}{\Delta}\right)$	$O\left(\frac{\log(T)}{\Delta}\right)$
Communication	0	$O\left(\log(T)\right)$	T

Even a minimal collaboration helps in reducing regret

Key Idea - Gossip the Best Arm

1. Agents use communication to recommend arms

Each message is log(K) bits

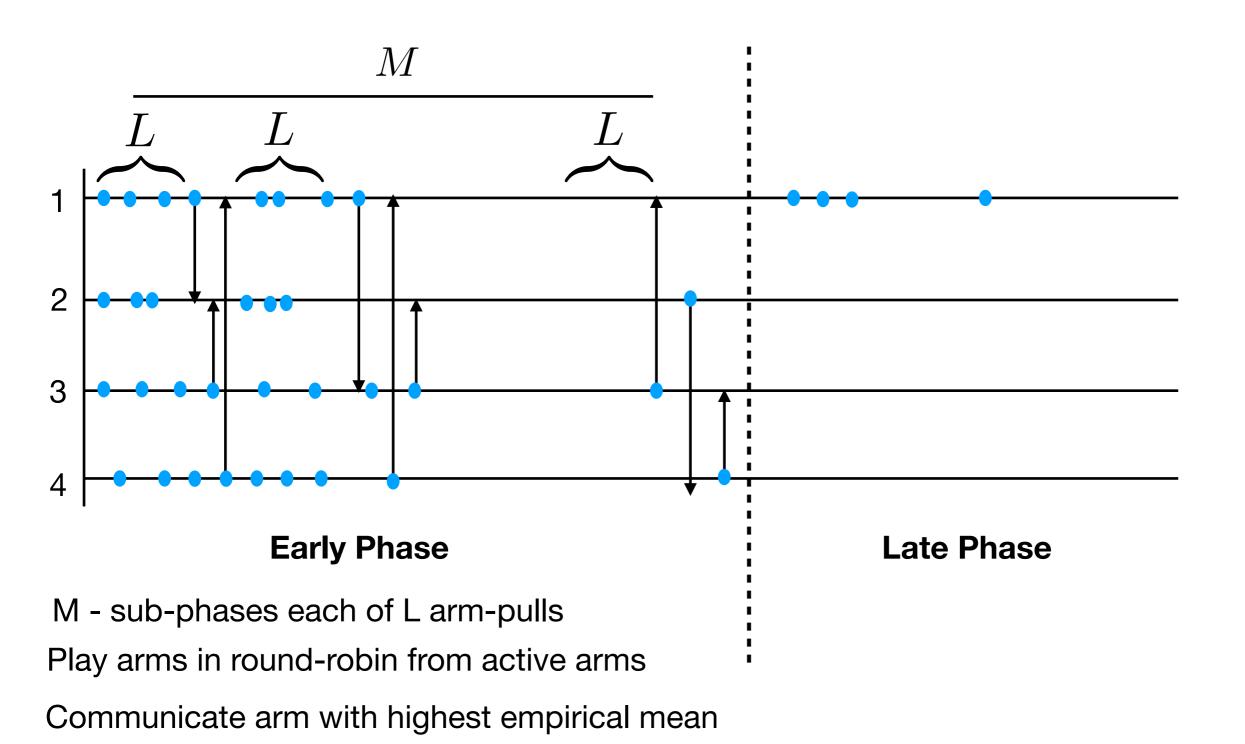
2. Agents only play from an active set of arms

Initially every agent has a small set of active arms that forms a <u>partition</u> of all arms (K/N)+1

The set of active arms grow with increasing recommendations (monotone)

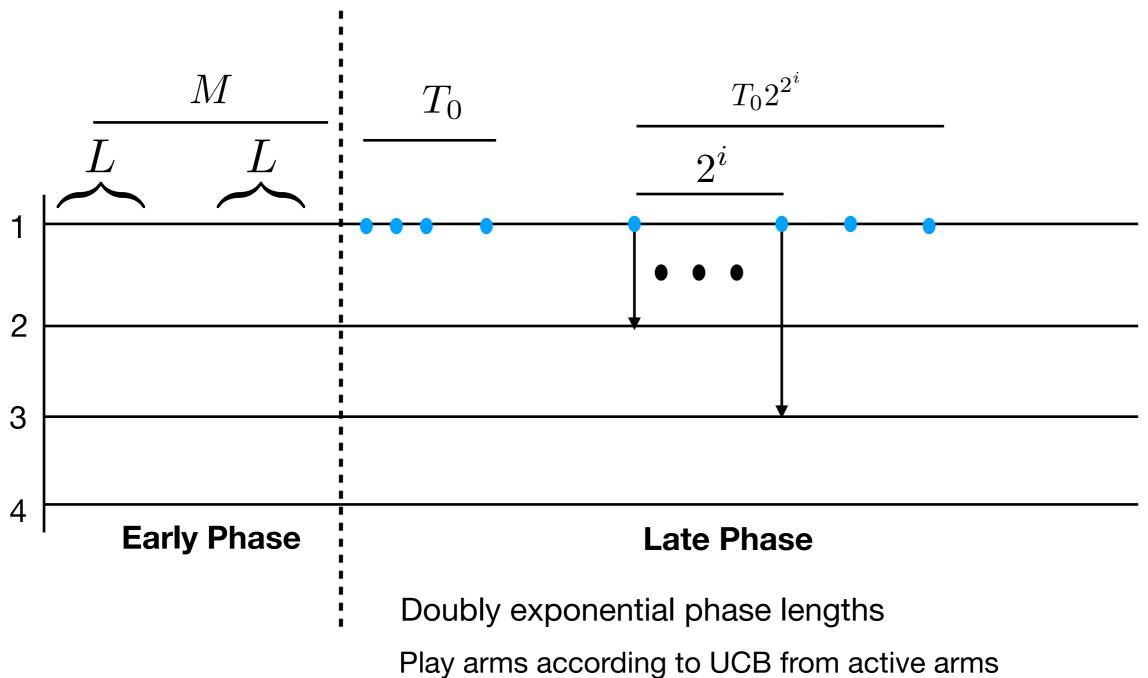
Algorithm Details

Two Phases to algorithm evolution



Algorithm Details

Two Phases to algorithm evolution



Communicate arm the most played arm in the previous phase

Performance Guarantees

Theorem

Suppose all agents execute the algorithm with parameters

$$M = O(\log(N))$$
, $0 < \varepsilon \le \Delta$, $L = O\left(\frac{\log^2(N)}{\varepsilon^2}\log\log(N)\right)$ and $T_0 = O\left(\frac{\max(K^2,N)}{\varepsilon^2}\log\left(\frac{1}{\varepsilon}\right)\right)$, then the regret of any agent i is

$$O\left(\left(\frac{K}{N} + \log(N)\right) \frac{\log(T)}{\Delta}\right) + \underbrace{O\left(\frac{\log^3(N)}{\varepsilon^2} \log\log(N)\right)}_{\text{Constant Independent of time}}$$

The total number of communications by any agent is $M + O(N \log(T))$

Performance Guarantees

Proof Ingredients

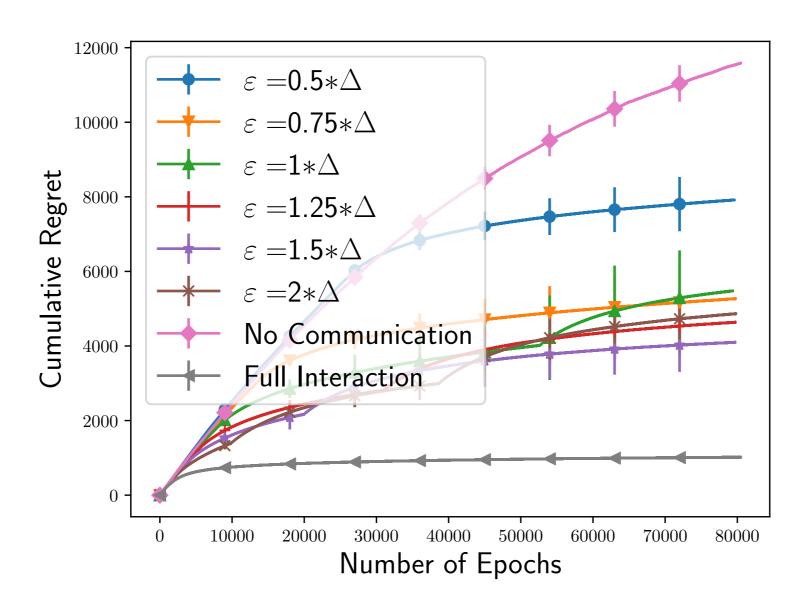
- 1. With high probability, at the end of the early phase, every agent's active arms
 - contains the best arm
 - total number of active arms is O(log(N) + K/N)
- 2. With high probability, at every late phase communication of every agent, the best arm is communicated.

<u>Algorithm Intuition</u>

Not every agent needs to play and figure out a bad arm

Amortize exploration cost of bad arms across the network

Empirical Performance



A representative plot with 20 agents and 50 Arms

Conclusions

A new collaborative multi-agent MAB setting

Social Learning based algorithm

Each agent plays from a reduced set of arms

Cost of exploration of a bad arm is amortized over the network