

Performance Oriented Association in Cellular Networks with Technology Diversity

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Outline

- Motivation and Background.
- Main mathematical model.
- Summary of Results.
- Conclusions

Background

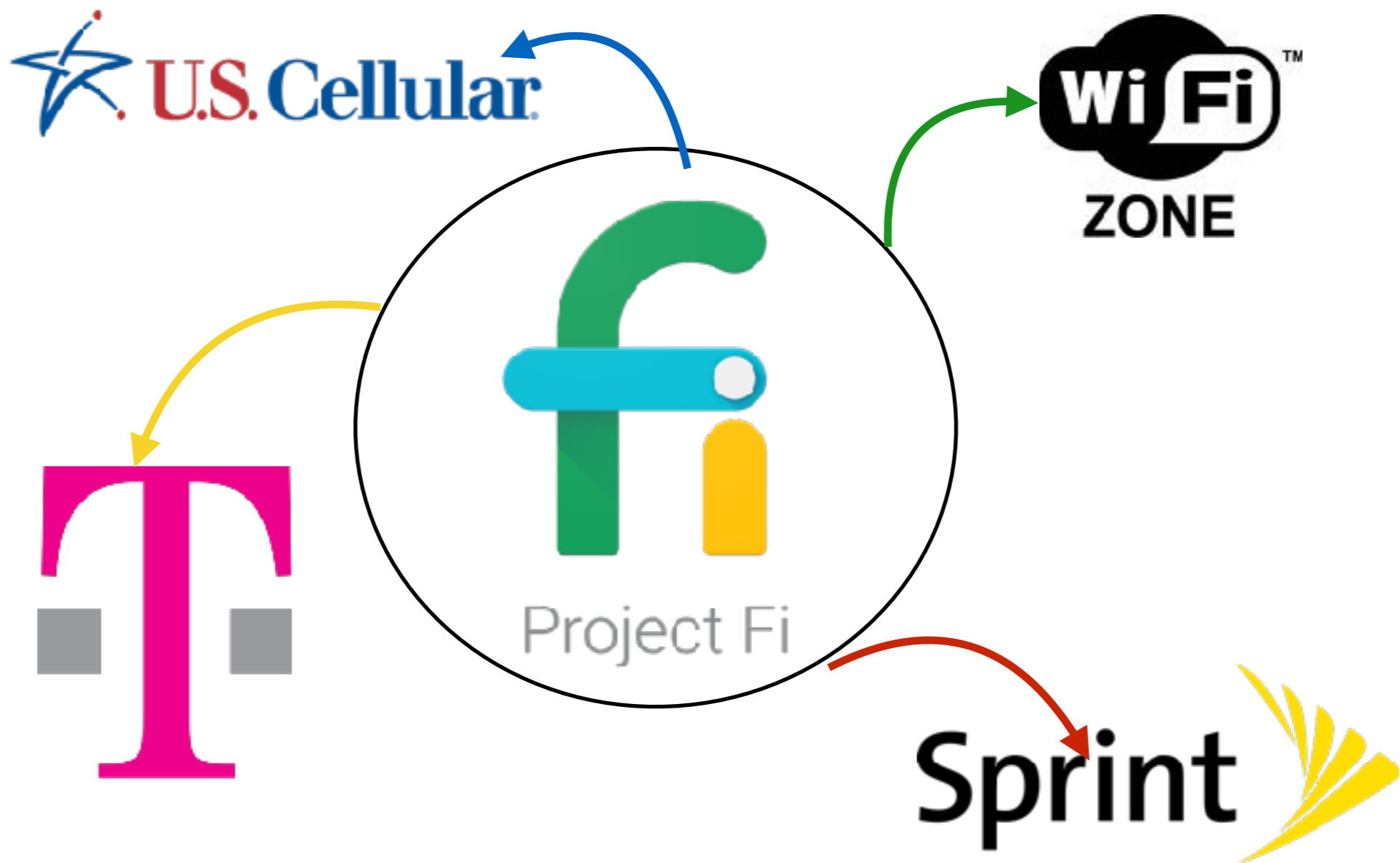
MVNOs (Mobile Virtual Network Operators)

These operators pool together various wireless technologies (for eg. 2G, 3G, 4G LTE) to create a service.

A UE can dynamically choose between the various technologies depending on whichever yields higher instantaneous benefit.



Background



A new model of cellular service by **Google**

What the Paper is About ? Technology Diversity

Leveraging the presence and *control* of multiple wireless technologies operating on *non-overlapping bandwidth*.

Google Fi - The different cellular operators and WiFi operate on separate bands.

MVNOs - Multiple orthogonal technologies (For ex. 3G and 4G LTE).

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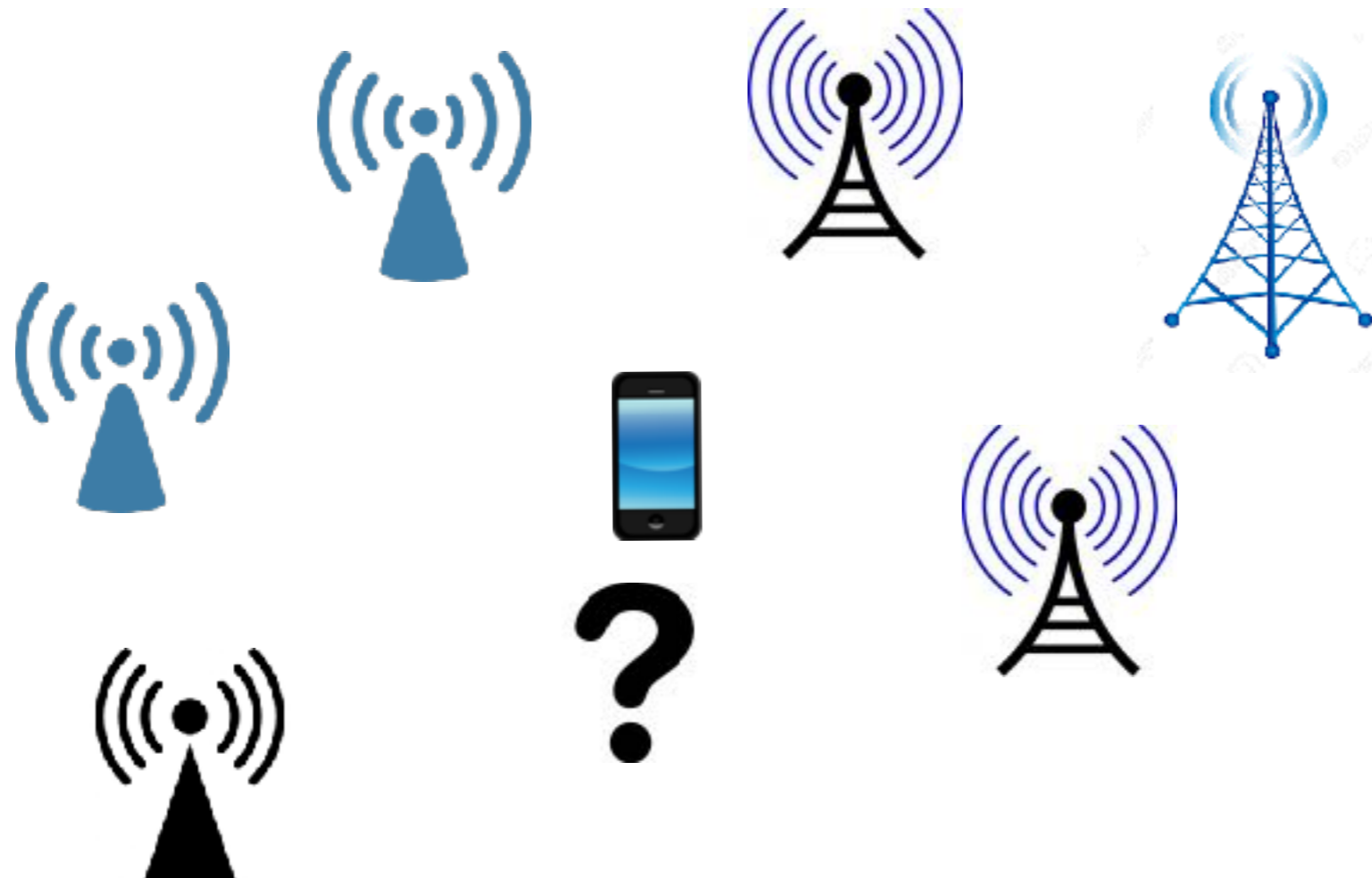
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Technology Diversity - *A framework to leverage and evaluate the benefits from the diversity of wireless technologies.*

The Problem we study - Base Station Association

Which Base-Station/Access Point must a UE associate with ?



A principled way to exploit the *diversity* in the network.

The “optimal” BS Association is not an obvious choice

First Guess - Connect to the nearest BS irrespective of the type of BS it is.

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Nearest BS



$$\text{SINR} = \frac{\text{Signal}}{\text{Interference} + \text{Noise}}$$

$$\text{SNR}_{\text{blue BS}} > \text{SNR}_{\text{black tower BS}}$$

(Basis for nearest BS association.)

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$$\text{SNR}_{\text{small BS}} > \text{SNR}_{\text{tower BS}} \quad (\text{Basis for nearest BS association.})$$

But bandwidths are not overlapping and thus interference only from one type of Base Stations.

$$\text{Thus, in this example } \text{SINR}_{\text{small BS}} < \text{SINR}_{\text{tower BS}}$$

Mathematical Framework - Network Model

Consider a network comprised of T different technologies.

The BS/APs of technology i is distributed as a independent Poisson Point Process $\phi_i \subset \mathbb{R}^2$ with intensity λ_i

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There is a *typical user* at the origin of the Euclidean plane who wishes to associate to a BS.

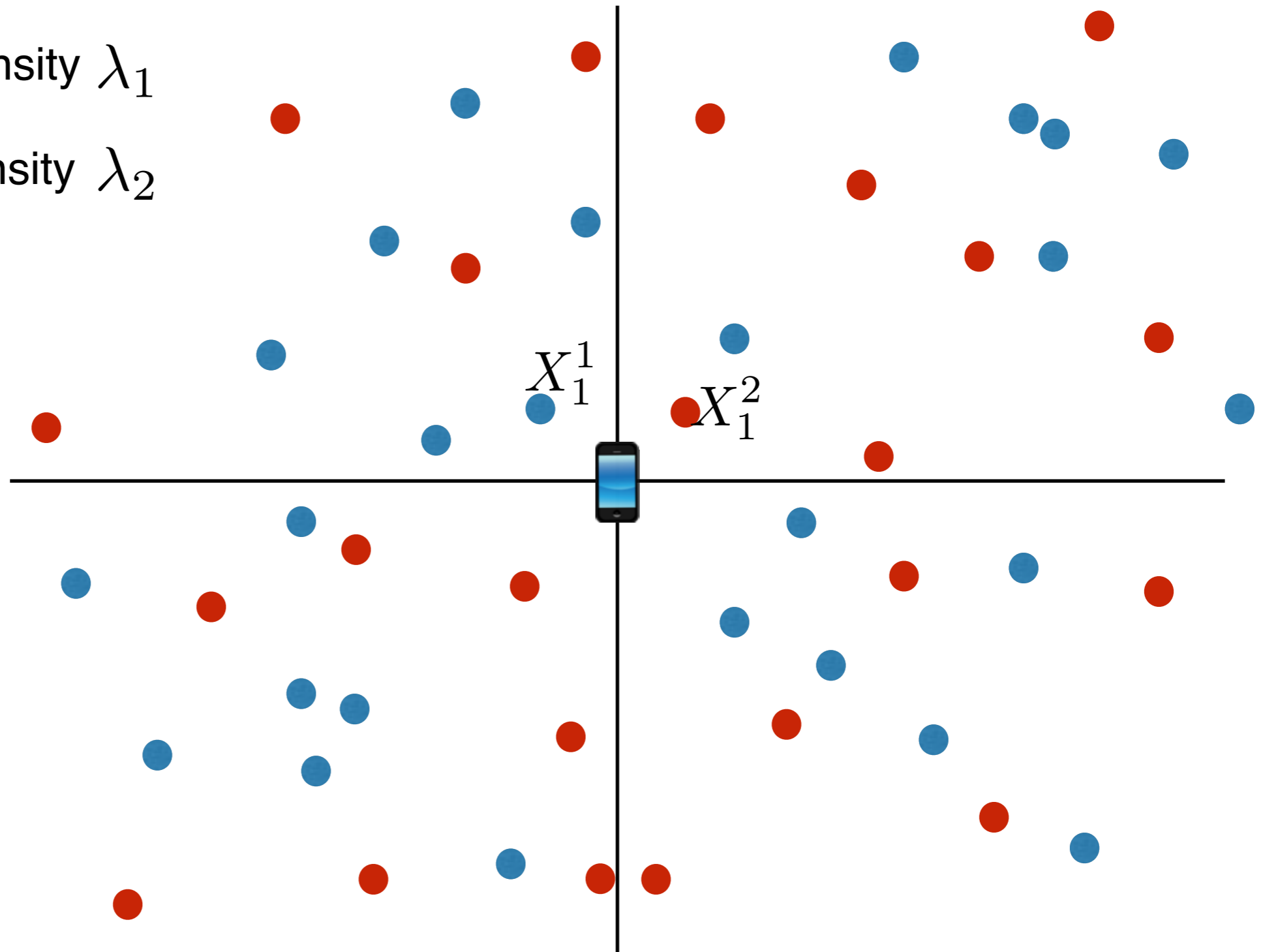
Palm theory connects the viewpoint of a single user to the **average** performance experienced by the users in the network.

Mathematical Framework - Network Model

$$T = 2$$

$\phi_1 = \bullet$ with intensity λ_1

$\phi_2 = \bullet$ with intensity λ_2



Mathematical Framework - Signal Model

BS of technology i transmits at power P_i

Signal from BS of technology i attenuated with distance as given by the function $l_i(\cdot) : \mathbb{R}_+ \rightarrow \mathbb{R}_+$

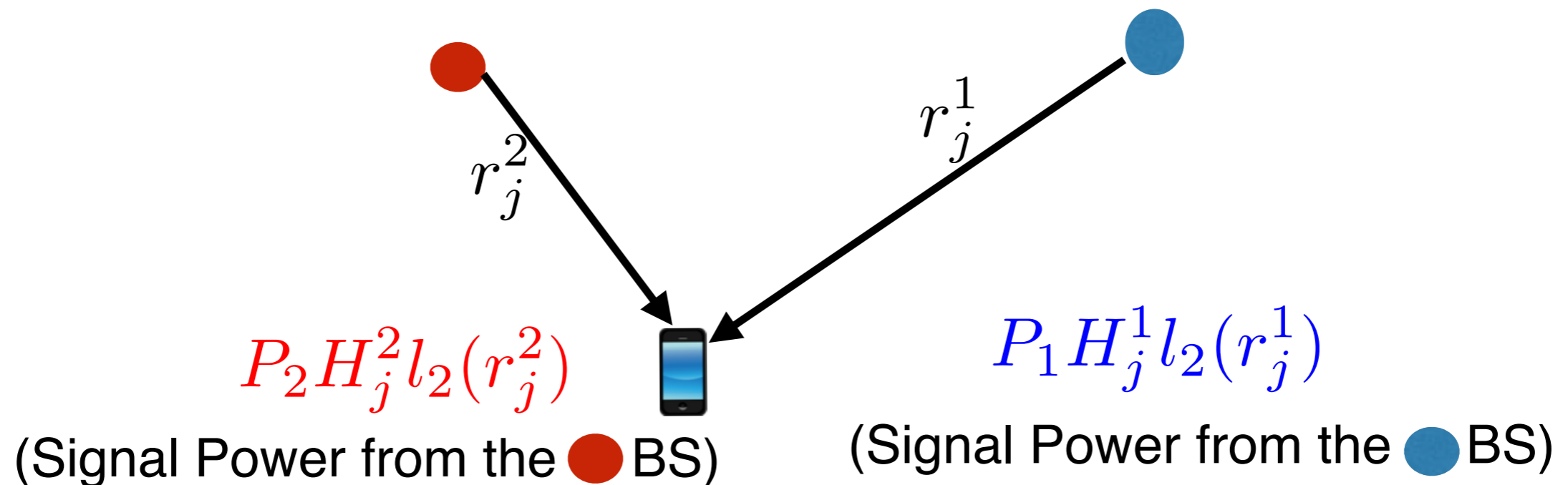
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Performance Metrics

(Non-overlapping bandwidths \implies Interference from only one technology)



$SINR_0^{i,j}$ SINR of the typical UE when it associates to the j th nearest BS of technology i .

- For each technology i , denote by bounded non-increasing functions $p_i(\cdot) : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ denoting the **reward** function.

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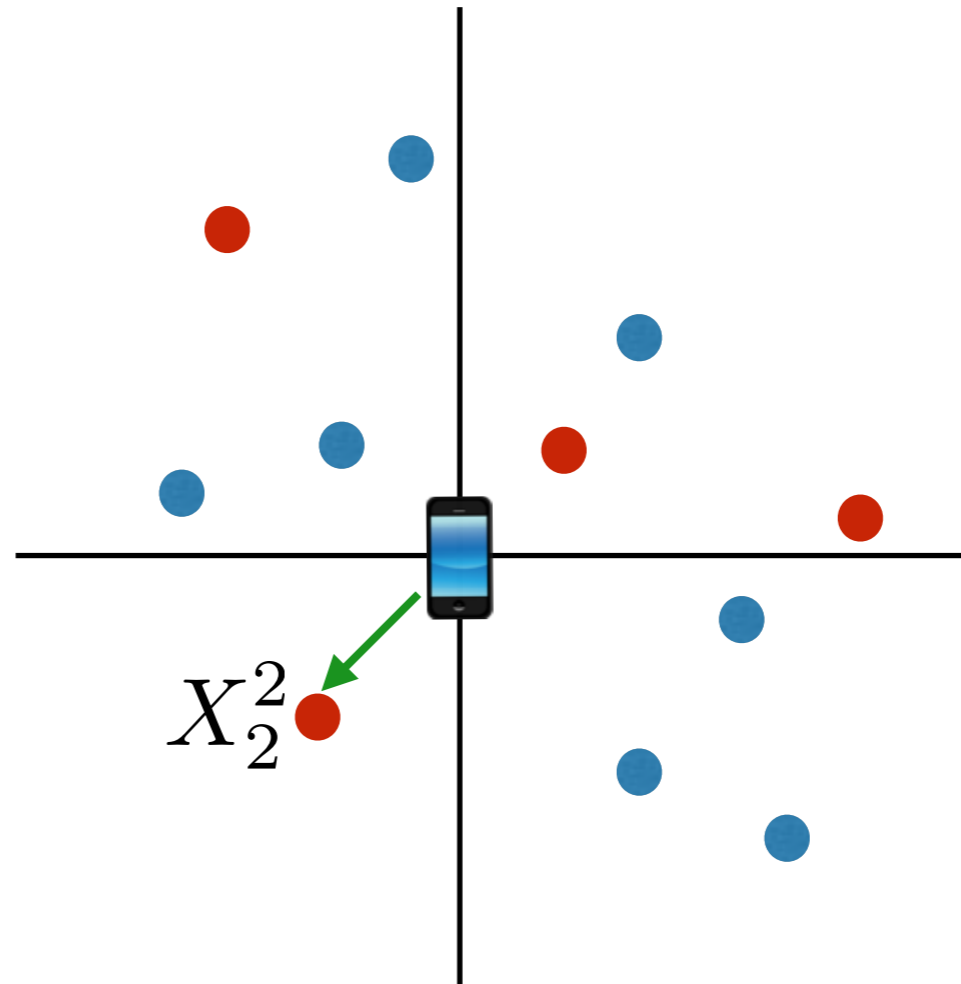
- For each technology i , denote by bounded non-increasing functions $p_i(\cdot) : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ denoting the **reward** function.
- If the typical UE connects to the j th nearest BS of technology i , then it receives a **reward** of $p_i(SINR_0^{i,j})$

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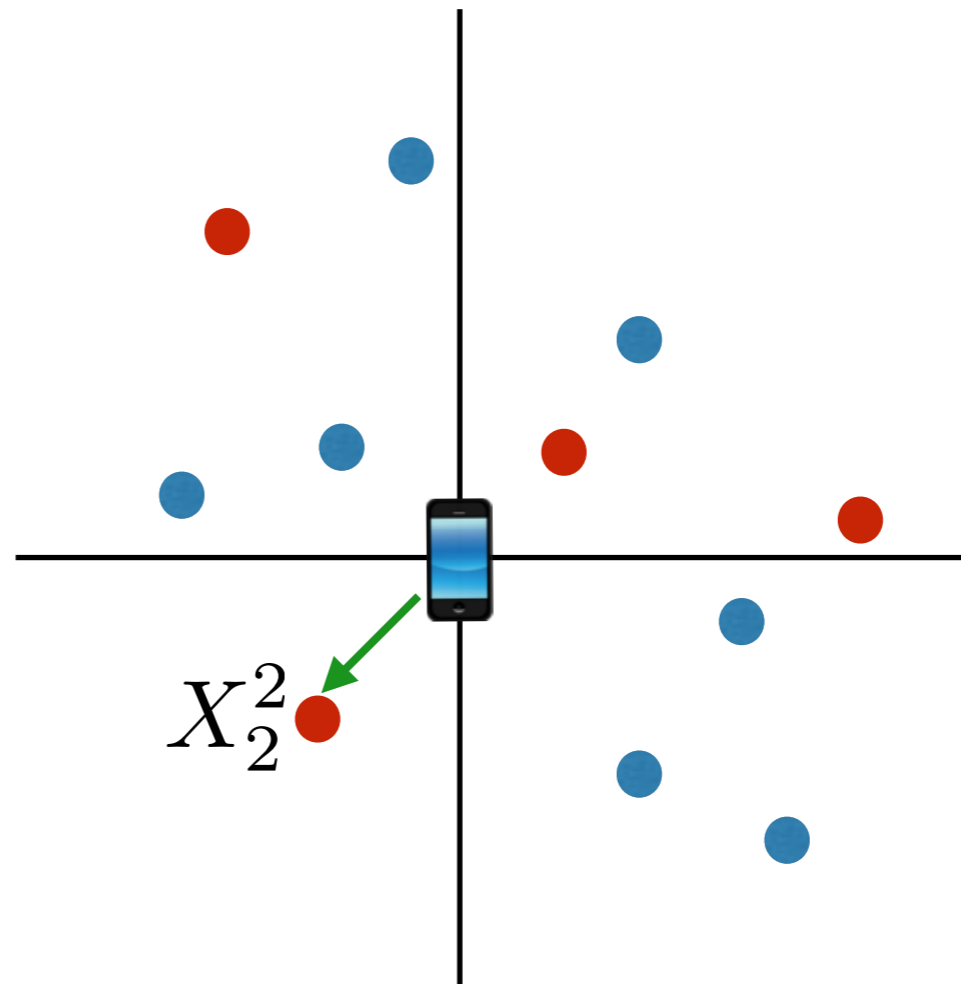
The reward received by the UE in this example is $p_2(\text{SINR}_0^{2,2})$

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Examples of common reward functions

- Coverage $p_i(x) = \mathbf{1}(x \geq \beta_i)$
- Average Achievable Rate $p_i(x) = B_i \log_2(1 + x)$

Information at the UE

Goal- Design association schemes exploiting available network *“information”* at the UE, that maximize expected reward of a typical UE.

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Examples of Information that a UE can know -

- Nearest BS of all technologies.
- Nearest k BS of all technologies.
- Instant fading and the distance to the nearest k BS
- Noisy estimate of the instant fading from the nearest k BS of each technology.

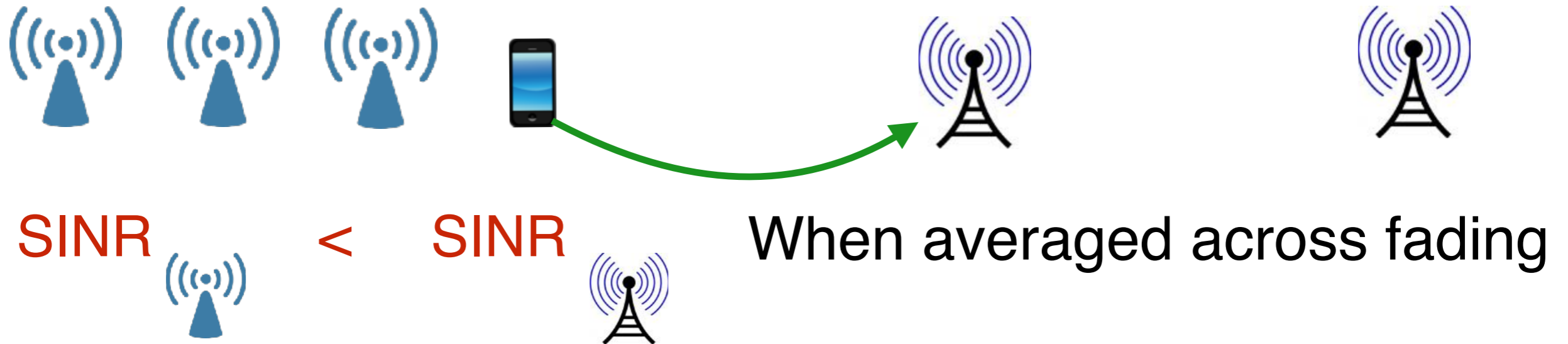
How Information affects Optimal Association



$$\text{SINR}_{\text{blue}} < \text{SINR}_{\text{black}}$$

When averaged across fading

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$SINR_{\text{small tower}} < SINR_{\text{large tower}}$ When averaged across fading



Sudden very good signal which the UE can sense.

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$SINR_{\text{nearby}} < SINR_{\text{distant}}$ When averaged across fading



Sudden very good signal which the UE can sense.

In this case, UE should associate to 

Information at the UE

Notion of Information at the typical UE formalized through the notion of *filtrations* of a sigma algebra.

Information at the UE - Formalization

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$(\Omega, \mathcal{F}, \mathbb{P})$ - The Probability space containing the independent RVs $\{\phi_i\}_{i=1}^T$ and $\{H_j^i\}_{i \in [1, T], j \in \mathbb{N}}$

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 $\mathcal{F}_I = \sigma \left(\sigma \left(\{r_k^i\}_{i=1}^T \right) \cup \mathcal{F}' \right)$ where $\mathcal{F}' \subseteq \sigma \left(\{H_j^i\}_{j \in [1, k], i \in [1, T]} \right)$

Performance Metric for Association

Recall -

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Performance of policy π is then $\mathcal{R}_\pi = \mathbb{E}[p_{\pi(0)}(SINR_0^\pi)]$

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Proposition:

$$\pi_I^* = \arg \sup_{i \in [1, T], j \in \mathbb{N}} \mathbb{E}[p_i(SINR_0^{i,j}) | \mathcal{F}_I] \quad \text{a.s.}$$

Optimal Association - Comparison of Information

The optimal association π_I^* under information \mathcal{F}_I is given by

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Comparison of schemes without messy computation !

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Optimal Association - Associate to the Nearest BS

Lemma:

If the information \mathcal{F}_I is independent of $\sigma(\{H_j^i\}_{j \geq 1, i \in [1, T]})$
then, $\arg \sup_{j \geq 1} \mathbb{E}[p_i(SINR_0^{i,j}) | \mathcal{F}_I] = 1$ for all i .

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*In particular, in the **absence of fading information**, the optimal strategy is to associate to the nearest BS of the best technology.*

Examples of Association Policies and Computations

Some examples of association policies.

- Optimal Association under no information.
- Optimal Association under nearest BS.
- Max-Ratio Association.
- Optimal Association under complete network information.

Max-Ratio Association

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$$\pi_m = \left[\arg \max_{i \in [1, T]} \frac{r_1^i}{r_2^i}, 1 \right]$$

Associate to the nearest BS of the technology yielding the maximum ratio.

(Oblivious to the statistical assumptions on the network.)

Max-Ratio Association

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Information - $\mathcal{F}_I = \sigma \left(\bigcup_{i=1}^T \frac{r_1^i}{r_2^i} \right)$

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This scheme trades off high signal power with that of interference power.



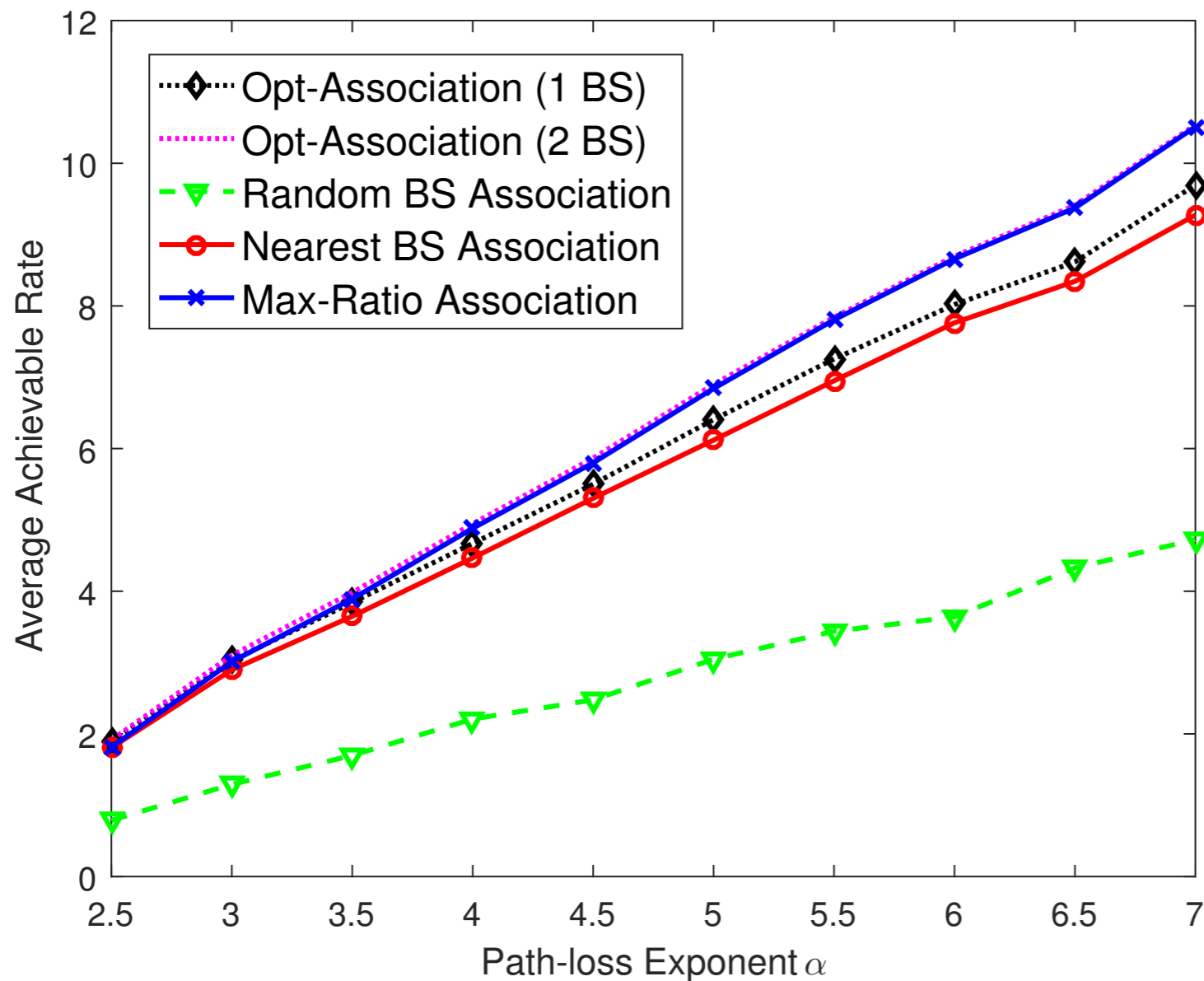
Max-Ratio Association - Asymptotic Optimality

Theorem - Let the noise powers be 0 and the reward function for all technologies are the same, i.e. $p_i(\cdot) = p(\cdot) \forall i$. Consider, the family of power law path-loss functions $\{l^{(\alpha)}(\cdot)\}_{\alpha > 2}$, where $l^{(\alpha)}(x) = x^{-\alpha}$. Let k be any integer larger than or equal to 2. If the information at the UE is the distance to the nearest k BS of each technology, i.e. $\mathcal{F}_I = \sigma(\cup_{i=1}^T (r_1^i, \dots, r_k^i))$, then Max-Ratio is asymptotically almost surely optimal i.e.

$$\pi_{\alpha}^* \xrightarrow{\alpha \rightarrow \infty} \left[\arg \max_{i \in [1, T]} \frac{r_1^i}{r_2^i}, 1 \right] \text{ a.s.}$$

Max-Ratio Association - Finite Scale Behavior

$$\pi_{\alpha}^* \xrightarrow{\alpha \rightarrow \infty} \left[\arg \max_{i \in [1, T]} \frac{r_1^i}{r_2^i}, 1 \right] \text{ a.s.}$$



For $\alpha \geq 5$, Max-Ratio is nearly optimal.

Formulas for Performance - Generalized Association

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Generalized Association - $i^* = \arg \max_{i \in [1, T]} \pi_i(\mathbf{r}_i, \lambda_i)$ and
by $j_i(\mathbf{r}_i, \lambda_i)$, the index of the BS to associate to if $i^* = i$.


Denote by $f_i(\cdot)$ as the PDF of the random variable $\pi_i(\mathbf{r}_i, \lambda_i)$
and by $F_i(\cdot)$ as the associated CDF.

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Denote by $f_i(\cdot)$ as the PDF of the random variable $\pi_i(\mathbf{r}_i, \lambda_i)$
and by $F_i(\cdot)$ as the associated CDF.


Randomness

Formulas for Performance - Generalized Association

Denote by \mathbf{r}_i , the L dimensional vector corresponding to the information about technology i

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Randomness

Example - The Max-Ratio and the Optimal Association fit into this performance evaluation framework.

Formulas for Performance - Generalized Association

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$f_i(\cdot), F_i(\cdot)$ The PDF and CDF of $\pi_i(\mathbf{r}_i, \lambda_i)$

Theorem -

$$\mathcal{R}_I^\pi = \sum_{i=1}^T \int_{\mathbf{r} \in \mathbb{R}^l} \mathbb{E}[p_i(\text{SINR}_0^{i, j_i(\mathbf{r}, \lambda_i)}) | \mathbf{r}] f_i(\mathbf{r}) \prod_{j=1, j \neq i}^T F_j(\pi_j(\mathbf{r}, \lambda_j)) d\mathbf{r}$$

Formulas for Performance - Generalized Association

\mathbf{r}_i The L dimensional vector corresponding to information on technology i

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Theorem -

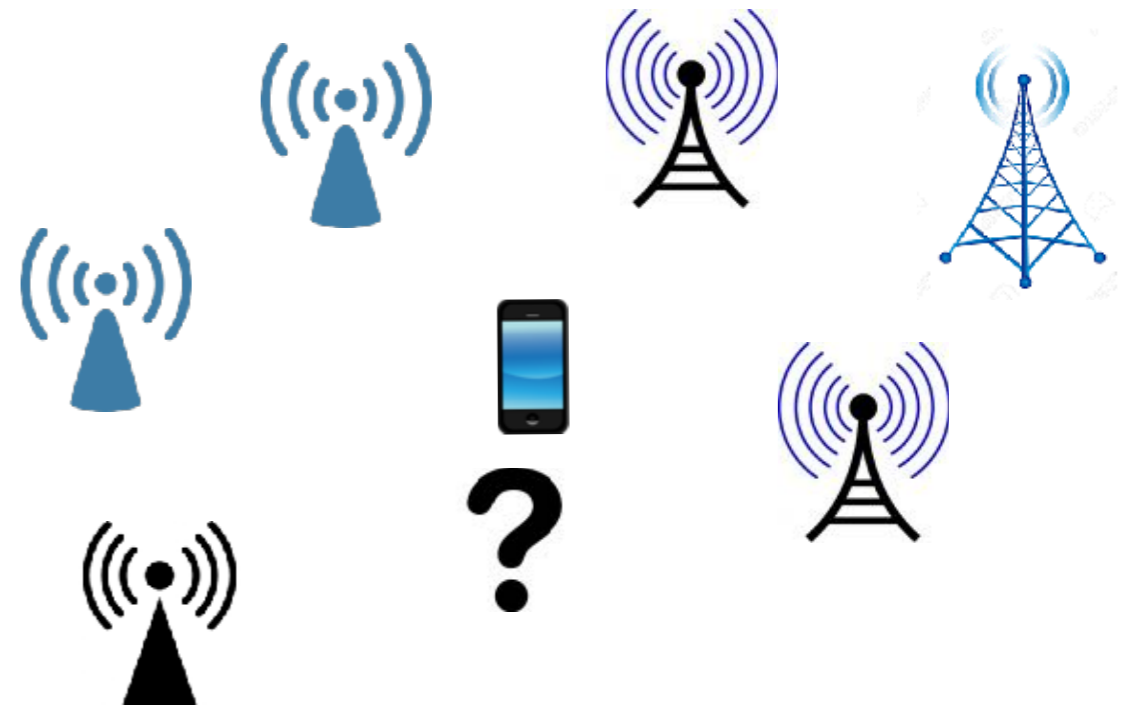
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Corollary -

Closed form formulas for Max-Ratio and certain optimal association schemes.

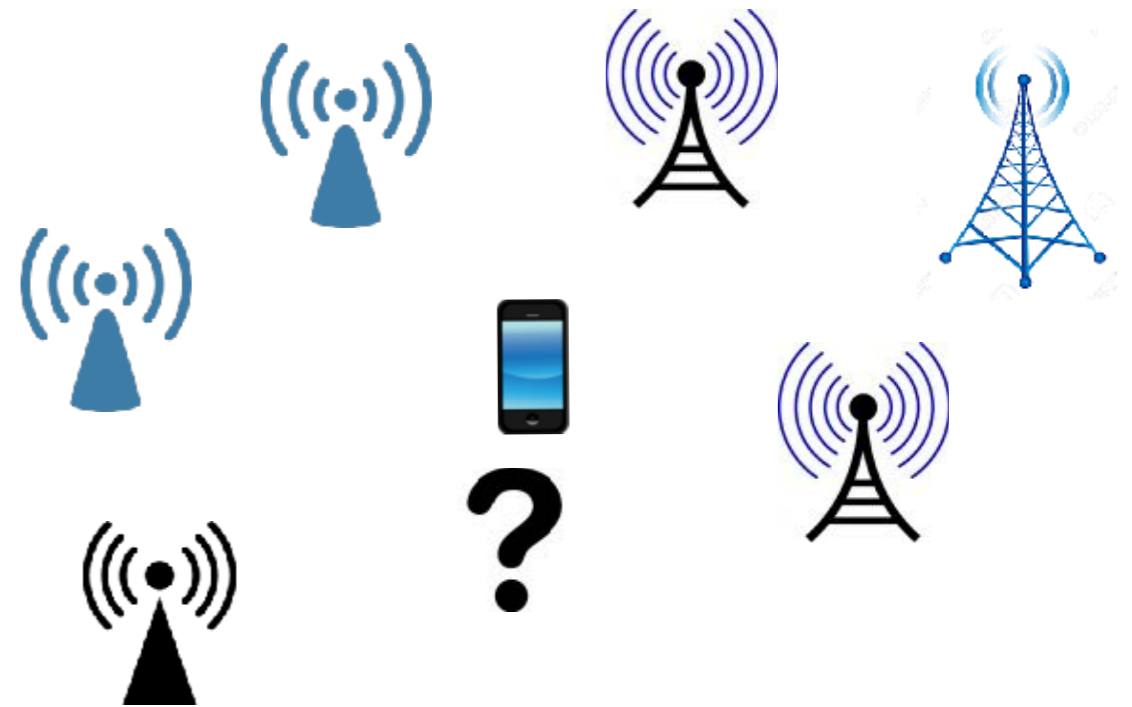
Conclusions

Provide a systematic framework to ***design*** and evaluate association schemes that exploits the available information and technology diversity



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Provide a systematic framework to ***design*** and evaluate association schemes that exploits the available information and technology diversity



Max-Ratio Algorithm -

A simple and *almost* optimal algorithm for association.

Thank you for your time.